One-shot signatures from noncollapsing hash functions **Based on joint work with Marcel Dall'Agnol**

Nick Spooner, 26th January 2024, University of Edinburgh

Collapsing hash functions*

$\sum_{\substack{\{0,1\}^n}} |x\rangle \longrightarrow h \longrightarrow \sum_{x \in \{0,1\}^n} |x, h(x)\rangle \longrightarrow \stackrel{\longrightarrow}{\underset{2nd}{\longrightarrow}} \sum_{x:h(x)=y}$ $x \in \{0, 1\}^n$

h is collapsing if blue and orange states are indistinguishable

Collapsing implies post-quantum collision resistance

*simplified definition



- Turns out to be the "correct" strengthening of CR for many applications (ZKPs, succinct arguments, etc.)

Big open question: does PQ-CR imply collapsing? = WIN > WIN =

- if yes, then collision resistance is enough for PQ security of protocols
 - (the answer is known to be yes for "regular" CRHFs [Zhandry22])
- if **no**, then we get cool quantum crypto!
- quantum lightning (strengthening of quantum money) [Zhandry19]
- one-shot signatures [DS23]



Equivocation game

Adversary

If **h** is collision-resistant, classical adversary wins w.p. <1/2+negl.

Can transform a quantum adversary winning w.p. >1/2+1/poly into OSS [DS23]

We show [DS23]: if h is not collapsing there is such an adversary



If *h* is collapsing, quantum adversary wins w.p. <1/2+negl. [Unruh16, CMSZ21]

Suppose *h* is **not** collapsing:

2nd $x \in \{0,1\}^n$ register





There is some index *i* such that:

$|x, h(x)\rangle$ $x \in \{0,1\}^n$ 2nd register















Applying orange measurement again gives **opposite** answer w.p. 1/poly!

Open question: are there non-collapsing PQCRHFs?

Many "natural" CRHFs are collapsing:

- Random functions [Unruh16]
- Ajtai hash function from LWE [LZ19, Poremba23]
- Polynomially-regular hash functions [Zhandry22]
- Optimally-secure CRHFs (PQ security $p(\lambda)/2^{\lambda}$) [Zhandry22]

- Potential route to constructing one-shot signatures (+ quantum money, etc.)

Non-collapsing CRHFs must be structured and sub-optimally secure.

